Category Theory & Programming

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ENTER FULLSCREEN

HTML presentation: use arrows, space to navigate.

Plan

- General overview

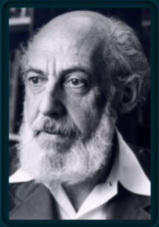
- Definitions

- Applications

General Overview

Recent Math Field

1942-45, Samuel Eilenberg & Saunders Mac Lane



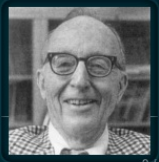
Certainly one of the more abstract branches of math

- *New math foundation*

  formalism abstraction, package entire theory[★]

- *Bridge between disciplines*

  Physics, Quantum Physics, Topology, Logic, Computer Science[☆]



★: When is one thing equal to some other thing?, Barry Mazur, 2007

☆: Physics, Topology, Logic and Computation: A Rosetta Stone, John C. Baez, Mike Stay, 2009

From a Programmer perspective

> *Category Theory is a new language/framework for Math*

Math Programming relation

Programming *is* doing Math

Not convinced?
Certainly a *vocabulary* problem.

One of the goal of Category Theory is to create a *homogeneous vocabulary* between different disciplines.

Vocabulary

Math vocabulary used in this presentation:

> *Category, Morphism, Associativity, Preorder, Functor, Endofunctor, Categorial property, Commutative diagram, Isomorph, Initial, Dual, Monoid, Natural transformation, Monad, Klesli arrows, κατα-morphism, ...*

Programmer Translation

Mathematician	Programmer
Morphism	Arrow
Monoid	String-like
Preorder	Acyclic graph
Isomorph	The same
Natural transformation	rearrangement function
Funny Category	LOLCat

Plan

- General overview
- Definitions
- Applications

- Category
- Intuition
- Examples
- Functor
- Examples

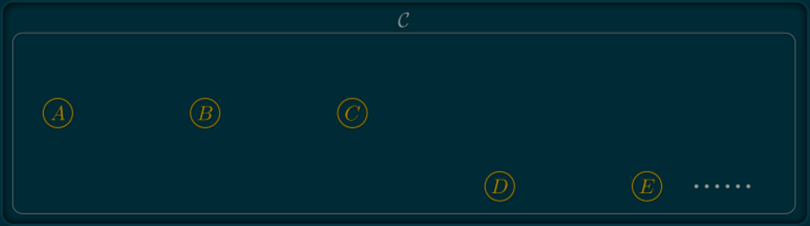
Category

A way of representing ***things*** and ***ways to go between things***.

A Category \(\mathcal{C}\) is defined by:

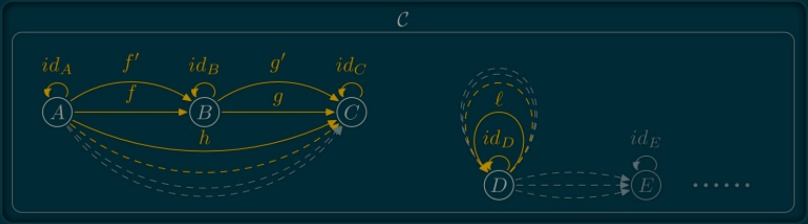
- *Objects \(\ob{C}\)*,
- *Morphisms \(\hom{C}\)*,
- a *Composition law* $(\circ)$
- obeying some *Properties*.

Category: Objects



\(\ob{\mathcal{C}}\) is a collection

Category: Morphisms



\(A\) and \(B\) objects of \(\C\)

\(\hom{A,B}\) is a collection of morphisms

\(f:A→B\) denote the fact \(f\) belongs to \(\hom{A,B}\)

\(\hom{\C}\) the collection of all morphisms of \(\C\)

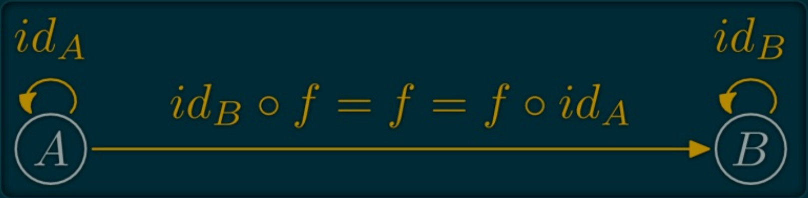
Category: Composition

Composition (∘): associate to each couple \(f:A→B, g:B→C\) $$g∘f:A\rightarrow C$$

$$\begin{tikzcd}
 & B \arrow[dr, bend left, "g"] & \\
A \arrow[ur, bend left, "f"] \arrow[rr, "g \circ f"] & & C
\end{tikzcd}$$

Category laws: neutral element

for each object \(X\), there is an \(\id_X:X→X\),
such that for each \(f:A→B\):



$$id_B \circ f = f = f \circ id_A$$

Category laws: Associativity

Composition is associative:

$$\begin{tikzcd}
A \arrow[r, "f"] \arrow[rr, bend left=40, "g \circ f"] \arrow[rrr, bend left=60, "(h \circ g) \circ f = h \circ (g \circ f)"] & B \arrow[r, "g"] \arrow[rr, bend right=40, "h \circ g"'] & C \arrow[r, "h"] & D
\end{tikzcd}$$

Commutative diagrams

Two path with the same source and destination are equal.

$$\begin{tikzcd}
A \arrow[r, "f"] \arrow[d, "g \circ f"'] & B \arrow[dl, "g"'] \arrow[d, "h \circ g"] \\
B \arrow[r, "h"] & C
\end{tikzcd}$$

\((h◦g)◦f = h◦(g◦f) \)

$$\begin{tikzcd}
A \arrow[r, "id_A"] \arrow[d, "f"'] & A \arrow[dl, "f"'] \arrow[d, "f"] \\
B \arrow[r, "id_B"] & B
\end{tikzcd}$$

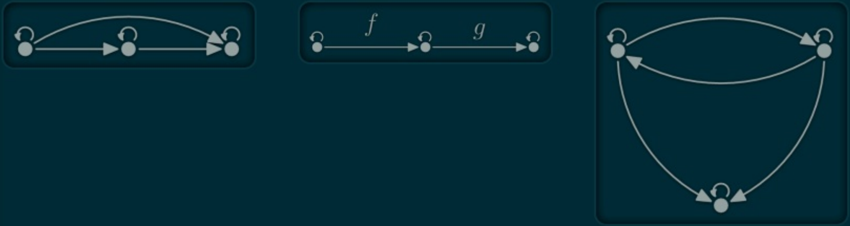
\(id_B◦f = f = f◦id_A \)

Question Time!

- French-only joke -

Can this be a category?

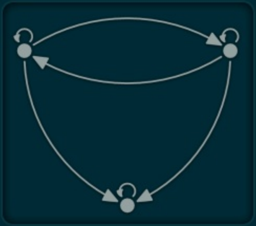
\(\ob{\C},\hom{\C}\) fixed, is there a valid $\circ$?



Can this be a category?

\(\ob{\C},\hom{\C}\) fixed, is there a valid $\circ$?

YES



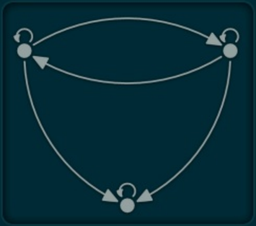
Can this be a category?

\(\ob{\C},\hom{\C}\) fixed, is there a valid ∘?

YES

no candidate for \(g∘f\)

NO



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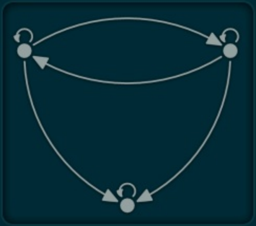
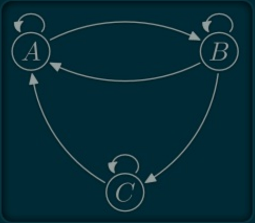
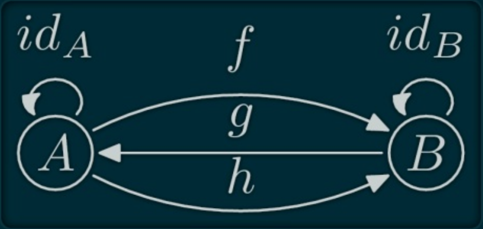


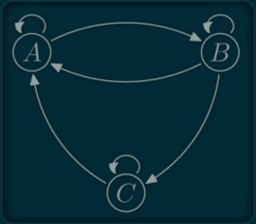
YES

Can this be a category?



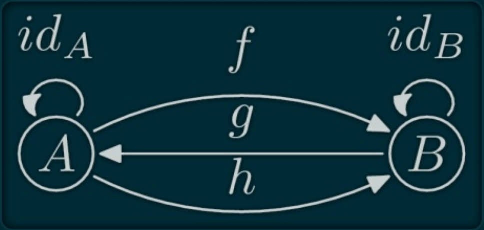


Can this be a category?

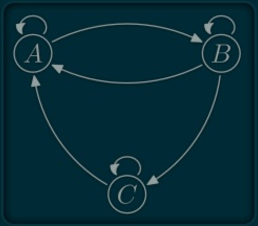


no candidate for \(f:C→B\)

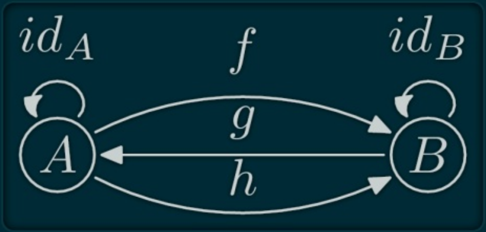
NO



Can this be a category?



no candidate for \(f:C→B\)
NO



\((h∘g)∘f=\id_B∘f=f\)
\(h∘(g∘f)=h∘\id_A=h\)
but \(h≠f\)
NO

Categories Examples

- Basket of Cats -

Category \(\Set\)

- \(\ob{\Set}\) are *all* the sets
- \(\hom{E,F}\) are *all* functions from \(E\) to \(F\)
- $\circ$ is functions composition

Category \(\Set\)

- \(\ob{\Set}\) are *all* the sets
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- $\circ$ is functions composition

- \(\ob{\Set}\) is a proper class ; not a set
- \(\hom{E,F}\) is a set
- \(\Set\) is then a *locally* ***small*** *category*

Categories Everywhere?

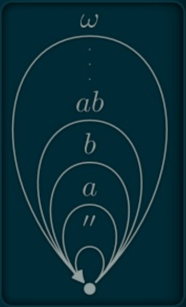
- \(\Mon\): (monoids, monoid morphisms, ∘)
- \(\Vec\): (Vectorial spaces, linear functions, ∘)
- \(\Grp\): (groups, group morphisms, ∘)
- \(\Rng\): (rings, ring morphisms, ∘)
- Any deductive system *T*: (theorems, proofs, proof concatenation)
- \(\Hask\): (Haskell types, functions, `(.)`)
- ...

Smaller Examples

Strings

- \(\ob{Str}\) is a singleton
- \(\hom{Str}\) each string
- ∘ is concatenation (++)

- "" ++ u = u = u ++ ""
- (u ++ v) ++ w = u ++ (v ++ w)

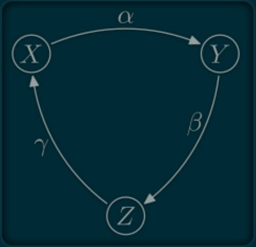


Finite Example?

Graph

- \(\ob{G}\) are vertices
- \(\hom{G}\) each path
- ∘ is path concatenation

- \(\ob{G}=\{X,Y,Z\}\),
- \(\hom{G}=\{ε,α,β,γ,αβ,βγ,...\}\)
- \(αβ∘γ=αβγ\)



Number construction

Each Numbers as a whole category

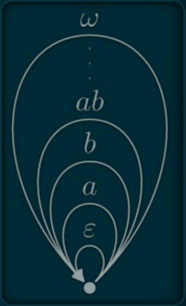
Degenerated Categories: Monoids

Each Monoid \((M,e,⊙): \ob{M}=\{ · \},\hom{M}=M,\circ = ⊙\)

Only one object.

Examples:

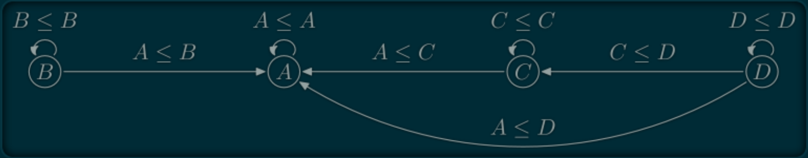
- `(Integer,0,+)`, `(Integer,1,*)`,
- `(Strings,"",++)`, for each `a`, `([a],[],++)`



Degenerated Categories: Preorders \((P,≤)\)

- \(\ob{P}={P}\),
- \(\hom{x,y}=\{x≤y\} ⇔ x≤y\),
- \((y≤z) \circ (x≤y) = (x≤z) \)

At most one morphism between two objects.



Degenerated Categories: Discrete Categories

Any Set

Any set \(E: \ob{E}=E, \hom{x,y}=\{x\} ⇔ x=y \)

Only identities

Categorical Properties

Any property which can be expressed in term of category, objects, morphism and composition.

- *Dual*: \(\D\) is \(\C\) with reversed morphisms.
- *Initial*: \(Z\in\ob{\C}\) s.t. \(∀Y∈\ob{\C}, \#\hom{Z,Y}=1\)
  Unique ("up to isormophism")
- *Terminal*: \(T\in\ob{\C}\) s.t. \(T\) is initial in the dual of \(\C\)
- *Functor*: structure preserving mapping between categories
- ...

Isomorph

isomorphism: $f:A\to B$ which can be "undone" *i.e.*
$\exists g:B\to A$, $g\circ f=id_A$ & $f\circ g=id_B$
in this case, A & B are *isomorphic*.

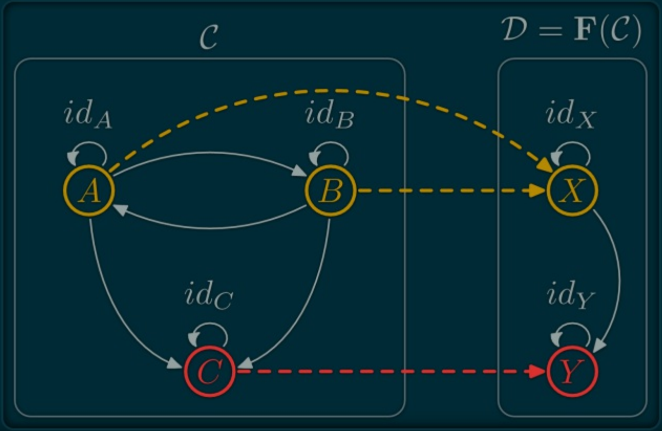
A≅B means A and B are essentially the same.
In Category Theory, = is in fact mostly ≅.
For example in commutative diagrams.

Functor

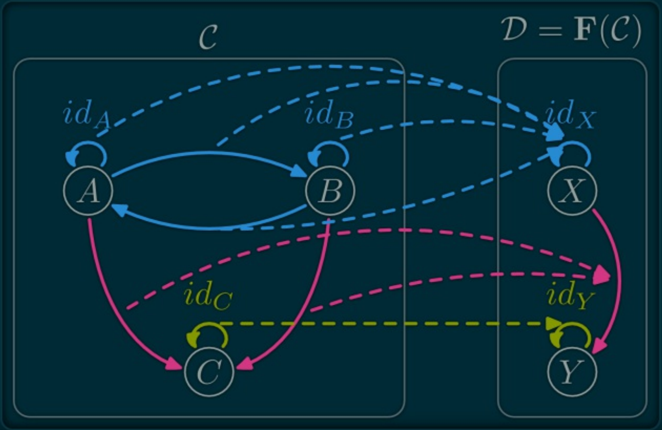
A functor is a mapping between two categories. Let \(\C\) and \(\D\) be two categories. A *functor* \(\F\) from \(\C\) to \(\D\):

- Associate objects: \(A\in\ob{\C}\) to \(\F(A)\in\ob{\D}\)
- Associate morphisms: \(f:A\to B\) to \(\F(f) : \F(A) \to \F(B)\) such that
    - \(\F (\)\(\id_X\)\()= \)\(\id\)\(\vphantom{\id}_{\F(}\)\(\vphantom{\id}_X\)\(\vphantom{\id}_{()} \),
    - \(\F (\)\(g∘f\)\()= \)\(\F(\)\(g\)\() \)\(\circ\)\(\F(\)\(f\)\() \)

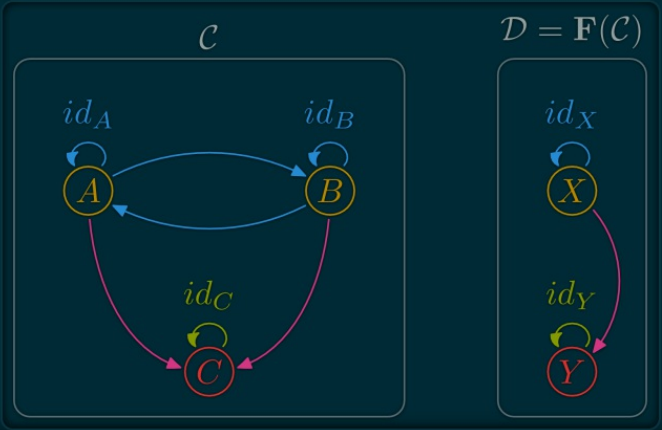
Functor Example (ob → ob)



Functor Example (hom → hom)

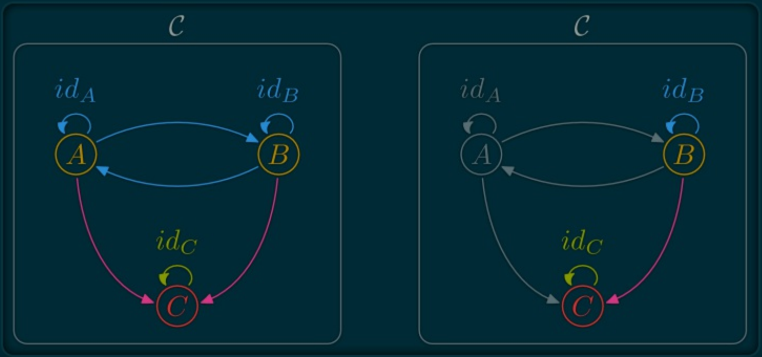


Functor Example



Endofunctors

An *endofunctor* for \(\C\) is a functor \(F:\C→\C\).



Category of Categories

Categories and functors form a category: \(\Cat\)

- \(\ob{\Cat}\) are categories
- \(\hom{\Cat}\) are functors
- ∘ is functor composition

Plan

- Why?
- What?
- **How?**

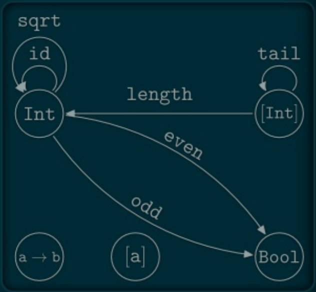
- \(\Hask\) category
- Functors
- Monads
- κατα-morphisms

Hask

Category \(\Hask\):

- \(\ob{\Hask} = \) Haskell types
- \(\hom{\Hask} = \) Haskell functions
- ∘ = `(.)` Haskell function composition

Forget glitches because of `undefined`.



Haskell Kinds

In Haskell some types can take type variable(s). Typically: `[a]`.

Types have *kinds*; The kind is to type what type is to function. Kind are the types for types (so meta).

```
Int, Char :: *
[], Maybe :: * -> *
(,) :: * -> * -> *
[Int], Maybe Char, Maybe [Int] :: *
```

Haskell Types

Sometimes, the type determine a lot about the function[★]:

```
fst :: (a,b) -> a -- Only one choice
snd :: (a,b) -> b -- Only one choice
f :: a -> [a]      -- Many choices
-- Possibilities: f x=[], or [x], or [x,x] or [x,...,x]

? :: [a] -> [a] -- Many choices
-- can only rearrange: duplicate/remove/reorder elements
-- for example: the type of addOne isn't [a] -> [a]
addOne l = map (+1) l
-- The (+1) force 'a' to be a Num.
```

[★]:Theorems for free!, Philip Wadler, 1989

Haskell Functor vs \(\Hask\) Functor

A Haskell Functor is a type `F :: * -> *` which belong to the type class `Functor` ; thus instantiate `fmap :: (a -> b) -> (F a -> F b)`.

`F`: \(\ob{\Hask}→\ob{\Hask}\)
& `fmap`: \(\hom{\Hask}→\hom{\Hask}\)

The couple `(F,fmap)` is a \(\Hask\)'s functor if for any `x :: F a`:

- `fmap id x = x`
- `fmap (f.g) x= (fmap f . fmap g) x`

Haskell Functors Example: Maybe

```
data Maybe a = Just a | Nothing
instance Functor Maybe where
    fmap :: (a -> b) -> (Maybe a -> Maybe b)
    fmap f (Just a) = Just (f a)
    fmap f Nothing = Nothing
```

```
fmap (+1) (Just 1) == Just 2
fmap (+1) Nothing  == Nothing
fmap head (Just [1,2,3]) == Just 1
```

Haskell Functors Example: List

```
instance Functor ([]) where
 fmap :: (a -> b) -> [a] -> [b]
 fmap = map
```

```
fmap (+1) [1,2,3]           == [2,3,4]
fmap (+1) []                == []
fmap head [[1,2,3],[4,5,6]] == [1,4]
```

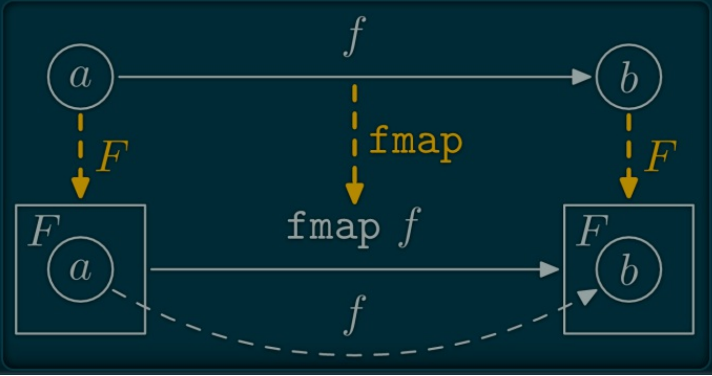
Haskell Functors for the programmer

`Functor` is a type class used for types that can be mapped over.

- Containers: `[]`, Trees, Map, HashMap...
- "Feature Type":
  - `Maybe a`: help to handle absence of `a`.
    Ex: `safeDiv x 0 ⇒ Nothing`
  - `Either String a`: help to handle errors
    Ex: `reportDiv x 0 ⇒ Left "Division by 0!"`

Haskell Functor intuition

Put normal function inside a container. Ex: list, trees...



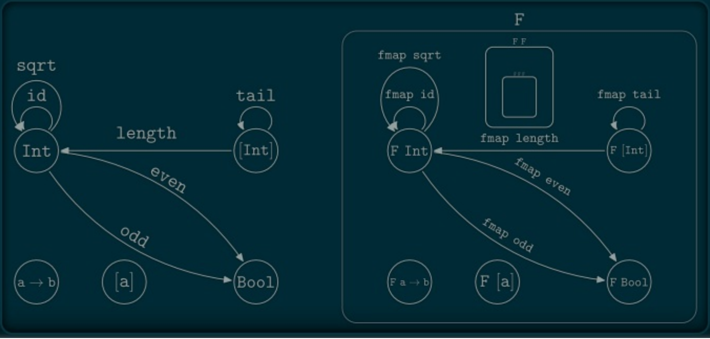
Haskell Functor properties

Haskell Functors are:

- *endofunctors* ; \(F:\C→\C\) here \(\C = \Hask\),
- a couple **(Object,Morphism)** in \(\Hask\).

Functor as boxes

Haskell functor can be seen as boxes containing all Haskell types and functions. Haskell types is fractal:



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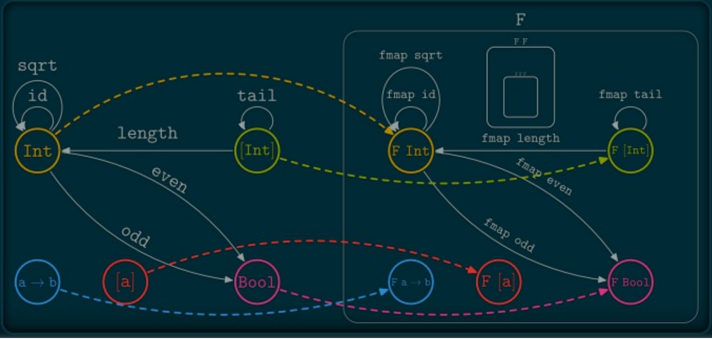

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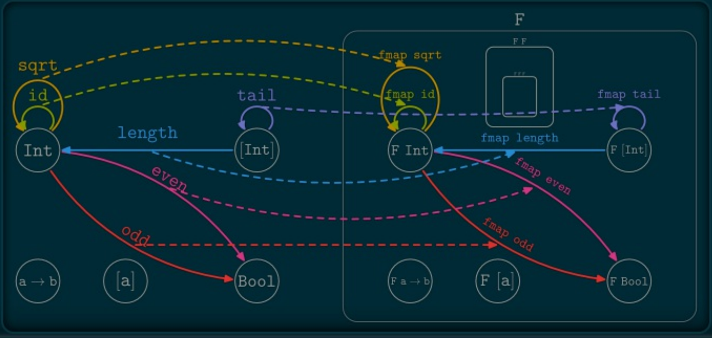


"Non Haskell" Hask's Functors

A simple basic example is the \(id_\Hask\) functor. It simply cannot be expressed as a couple (`F`,`fmap`) where

- `F::* -> *`
- `fmap :: (a -> b) -> (F a) -> (F b)`

Another example:

- F(`T`)=`Int`
- F(`f`)=`\_->0`

Also Functor inside \(\Hask\)

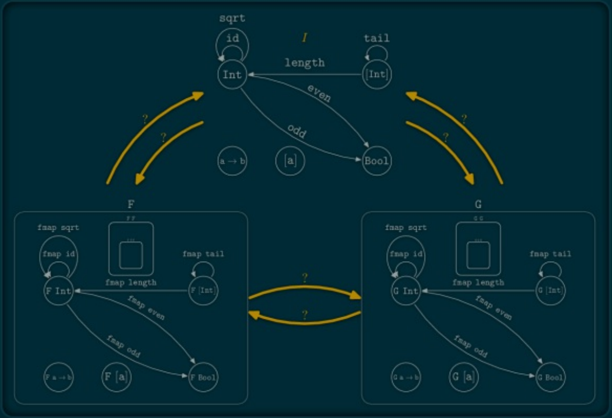
\(\mathtt{[a]}∈\ob{\Hask}\) but is also a category. Idem for `Int`.

`length` is a Functor from the category `[a]` to the cateogry `Int`:

- \(\ob{\mathtt{[a]}}=\{ · \}\)		- \(\ob{\mathtt{Int}}=\{ · \}\)
- \(\hom{\mathtt{[a]}}=\mathtt{[a]}\)	⇒	- \(\hom{\mathtt{Int}}=\mathtt{Int}\)
- \(∘=\mathtt{(++)}\)		- \(∘=\mathtt{(+)}\)

- id: `length [] = 0`
- comp: `length (l ++ l') = (length l) + (length l')`

Category of \(\Hask\) Endofunctors



Category of Functors

If \(\C\) is *small* (\(\hom{\C}\) is a set). All functors from \(\C\) to some category \(\D\) form the category \(\mathrm{Func}(\C,\D)\).

- \(\ob{\mathrm{Func}(\C,\D)}\): Functors \(F:\C→\D\)
- \(\hom{\mathrm{Func}(\C,\D)}\): *natural transformations*
- ∘: Functor composition

\(\mathrm{Func}(\C,\C)\) is the category of endofunctors of \(\C\).

Natural Transformations

Let \(F\) and \(G\) be two functors from \(\C\) to \(\D\).

A natural transformation: familly η ; \(η_X\in\hom{\D}\) for \(X\in\ob{\C}\) s.t.

$$\begin{tikzcd}
FX \arrow[r, "Ff"] \arrow[d, "\eta_X"'] & FY \arrow[d, "\eta_Y"] \\
GX \arrow[r, "Gf"] & GY
\end{tikzcd}$$

ex: between Haskell functors; `F a -> G a`
Rearragement functions only.

Natural Transformation Examples (1/4)

```
data Tree a = Empty | Node a [Tree a]
toTree :: [a] -> Tree a
toTree [] = Empty
toTree (x:xs) = Node x [toTree xs]
```

`toTree` is a natural transformation. It is also a morphism from `[]` to `Tree` in the Category of \(\Hask\) endofunctors.

$$\begin{tikzcd}
{[a]} \arrow[r, "\texttt{fmap}_{[]}\ \texttt{f}"] \arrow[d, "\texttt{toTree}"'] & {[b]} \arrow[d, "\texttt{toTree}"] \\
\texttt{Tree a} \arrow[r, "\texttt{fmap}_{\texttt{Tree}}\ \texttt{f}"] & \texttt{Tree b}
\end{tikzcd}$$

Natural Transformation Examples (2/4)

```
data Tree a = Empty | Node a [Tree a]
toList :: Tree a -> [a]
toList Empty = []
toList (Node x l) = [x] ++ concat (map toList l)
```

`toList` is a natural transformation. It is also a morphism from `Tree` to `[]` in the Category of \(\Hask\) endofunctors.

$$\begin{tikzcd}
{[a]} \arrow[r, "\mathtt{fmap}_{[]}\ \mathtt{f}"] \arrow[d, shift right, "\mathtt{toTree}"'] & {[b]} \arrow[d, shift left, "\mathtt{toTree}"] \\
\mathtt{Tree\ a} \arrow[r, "\mathtt{fmap}_{\mathtt{Tree}}\ \mathtt{f}"] \arrow[u, shift right, "\mathtt{toList}"'] & \mathtt{Tree\ b} \arrow[u, shift left, "\mathtt{toList}"]
\end{tikzcd}$$

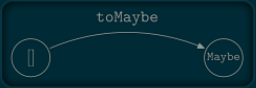
toList . toTree = id & toTree . toList = id
therefore [] & Tree are isomorph.

Natural Transformation Examples (3/4)

```
toMaybe :: [a] -> Maybe a
toMaybe [] = Nothing
toMaybe (x:xs) = Just x
```

`toMaybe` is a natural transformation. It is also a morphism from `[]` to `Maybe` in the Category of \(\Hask\) endofunctors.

$$\begin{tikzcd}
{[a]} \arrow[r, "\mathtt{fmap}_{[]}\ \mathtt{f}"] \arrow[d, "\mathtt{toMaybe}"'] & {[b]} \arrow[d, "\mathtt{toMaybe}"] \\
\mathtt{Maybe\ a} \arrow[r, "\mathtt{fmap}_{\mathtt{Maybe}}\ \mathtt{f}"] & \mathtt{Maybe\ b}
\end{tikzcd}$$

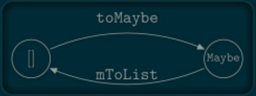


Natural Transformation Examples (4/4)

```
mToList :: Maybe a -> [a]
mToList Nothing = []
mToList Just x  = [x]
```

`toMaybe` is a natural transformation. It is also a morphism from `[]` to `Maybe` in the Category of \(\Hask\) endofunctors.

$$\begin{tikzcd}
{[a]} \arrow[r, "\mathtt{fmap}_{[]}\ \mathtt{f}"] & {[b]} \\
\mathtt{Maybe\ a} \arrow[u, "\mathtt{mToList}"] \arrow[r, "\mathtt{fmap}_{\mathtt{Maybe}}\ \mathtt{f}"] & \mathtt{Maybe\ b} \arrow[u, "\mathtt{mToList}"']
\end{tikzcd}$$



There is no isomorphism.
Hint: Bool lists longer than 1.

Composition problem

The Problem; example with lists:

```
f x = [x]         ⇒ f 1 = [1]   ⇒ (f.f) 1 = [[1]] ✗
g x = [x+1]       ⇒ g 1 = [2]   ⇒ (g.g) 1 = ERROR [2]+1 ✗
h x = [x+1,x*3] ⇒ h 1 = [2,3] ⇒ (h.h) 1 = ERROR [2,3]+1 ✗
```

The same problem with most `f :: a -> F a` functions and functor `F`.

Composition Fixable?

How to fix that? We want to construct an operator which is able to compose:

`f :: a -> F b` *&* `g :: b -> F c`.

More specifically we want to create an operator ◎ of type

`◎ :: (b -> F c) -> (a -> F b) -> (a -> F c)`

Note: if `F` = I, ◎ = `(.)`.

Fix Composition (1/2)

Goal, find: `◎ :: (b -> F c) -> (a -> F b) -> (a -> F c)`

`f :: a -> F b`, `g :: b -> F c`:

- `(g ◎ f) x` ???
- First apply `f` to `x` ⇒ `f x :: F b`
- Then how to apply `g` properly to an element of type `F b`?

Fix Composition (2/2)

Goal, find: `◎ :: (b -> F c) -> (a -> F b) -> (a -> F c)`

`f :: a -> F b`, `g :: b -> F c`, `f x :: F b`:

- Use `fmap :: (t -> u) -> (F t -> F u)`!
- `(fmap g) :: F b -> F (F c)` ; (`t=b`, `u=F c`)
- `(fmap g) (f x) :: F (F c)` it almost WORKS!
- We lack an important component, `join :: F (F c) -> F c`
- `(g ◎ f) x = join ((fmap g) (f x))` ☺

  ◎ is the Kleisli composition; in Haskell: `<=<` (in `Control.Monad`).

Necessary laws

For ◎ to work like composition, we need join to hold the following properties:

- `join (join (F (F (F a))))=join (F (join (F (F a))))`
- abusing notations denoting `join` by ⊙; this is equivalent to
  $(F \odot F) \odot F = F \odot (F \odot F)$
- There exists `η :: a -> F a` s.t.
  $\eta \odot F = F = F \odot \eta$

Klesli composition

Now the composition works as expected. In Haskell ◎ is `<=<` in `Control.Monad`.

```
g <=< f = \x -> join ((fmap g) (f x))
```

```
f x = [x]       ⇒ f 1 = [1]   ⇒ (f <=< f) 1 = [1] ✓
g x = [x+1]     ⇒ g 1 = [2]   ⇒ (g <=< g) 1 = [3] ✓
h x = [x+1,x*3] ⇒ h 1 = [2,3] ⇒ (h <=< h) 1 = [3,6,4,9] ✓
```

We reinvented Monads!

A monad is a triplet `(M,⊙,η)` where

- \(M\) an Endofunctor (to type `a` associate `M a`)
- \(⊙:M×M→M\) a nat. trans. (i.e. `⊙::M (M a) → M a` ; `join`)
- \(η:I→M\) a nat. trans. (\(I\) identity functor ; `η::a → M a`)

Satisfying

- \(M ⊙ (M ⊙ M) = (M ⊙ M) ⊙ M\)
- \(η ⊙ M = M = M ⊙ η\)

Compare with Monoid

A Monoid is a triplet \((E, \cdot ,e)\) s.t.

- \(E\) a set
- \(\cdot : E \times E \to E\)
- \(e : 1 \to E\)

Satisfying

- \(x \cdot (y \cdot z) = (x \cdot y) \cdot z, \ \forall x,y,z \in E\)
- \(e \cdot x = x = x \cdot e, \ \forall x \in E\)

Monads are just Monoids

> *A Monad is just a monoid in the category of endofunctors, what's the problem?*

The real sentence was:

> *All told, a monad in X is just a monoid in the category of endofunctors of X, with product × replaced by composition of endofunctors and unit set by the identity endofunctor.*

Example: List

- `[] :: * -> *` an Endofunctor
- \(⊙:M×M→M\) a nat. trans. (`join :: M (M a) -> M a`)
- \(η:I→M\) a nat. trans.

```
-- In Haskell ⊙ is "join" in "Control.Monad"
join :: [[a]] -> [a]
join = concat

-- In Haskell the "return" function (unfortunate name)
η :: a -> [a]
η x = [x]
```

Example: List (law verification)

Example: `List` is a functor (`join` is ⊙)

- $M \odot (M \odot M) = (M \odot M) \odot M$
- $\eta \odot M = M = M \odot \eta$

```
join [ join [[x,y,...,z]] ] = join [[x,y,...,z]]
                      = join (join [[[x,y,...,z]]])
join (η [x]) = [x] = join [η x]
```

Therefore `([],join,η)` is a monad.

Monads useful?

A *LOT* of monad tutorial on the net. Just one example; the State Monad

`DrawScene` to `State Screen DrawScene` ; still **pure**.

```
main = drawImage (width,height)

drawImage :: Screen -> DrawScene
drawImage screen =
    drawPoint p screen
    drawCircle c screen
    drawRectangle r screen

drawPoint point screen = ...
drawCircle circle screen = ...
drawRectangle rectangle screen = ...
```

```
main = do
    put (Screen 1024 768)
    drawImage

drawImage :: State Screen DrawScene
drawImage = do
    drawPoint p
    drawCircle c
    drawRectangle r

drawPoint :: Point -> State Screen DrawScene
drawPoint p = do
    Screen width height <- get
    ...
```

fold

κατα-morphism

κατα-morphism: fold generalization

`acc` type of the "accumulator":

`fold :: (acc -> a -> acc) -> acc -> [a] -> acc`

Idea: put the accumulated value inside the type.

```
-- Equivalent to fold (+1) 0 "cata"
(Cons 'c' (Cons 'a' (Cons 't' (Cons 'a' Nil))))
(Cons 'c' (Cons 'a' (Cons 't' (Cons 'a' 0))))
(Cons 'c' (Cons 'a' (Cons 't' 1)))
(Cons 'c' (Cons 'a' 2))
(Cons 'c' 3)
4
```

But where are all the informations? `(+1)` and `0`?

κατα-morphism: Missing Information

Where is the missing information?

- Functor operator `fmap`
- Algebra representing the `(+1)` and also knowing the `0`.

First example, make `length` on `[Char]`

ката-morphism: Type work

```
data StrF a = Cons Char a | Nil
data Str = StrF Str

-- generalize the construction of Str to other datatype
-- Mu :: type fixed point

data Mu f = InF { outF :: f (Mu f) }
data Str = Mu StrF

-- Example
foo=InF { outF = Cons 'f'
        (InF { outF = Cons 'o'
            (InF { outF = Cons 'o'
                (InF { outF = Nil })})})}
```

κατα-morphism: missing information retrieved

```
type Algebra f a = f a -> a
instance Functor (StrF a) =
    fmap f (Cons c x) = Cons c (f x)
    fmap _ Nil = Nil
```

```
cata :: Functor f => Algebra f a -> Mu f -> a
cata f = f . fmap (cata f) . outF
```

κατα-morphism: Finally length

All needed information for making length.

```haskell
instance Functor (StrF a) =
    fmap f (Cons c x) = Cons c (f x)
    fmap _ Nil = Nil

length' :: Str -> Int
length' = cata phi where
    phi :: Algebra StrF Int -- StrF Int -> Int
    phi (Cons a b) = 1 + b
    phi Nil = 0

main = do
    l <- length' $ stringToStr "Toto"
    ...
```

#ката-morphism: extension to Trees

Once you get the trick, it is easy to extent to most Functor.

```
type Tree = Mu TreeF
data TreeF x = Node Int [x]

instance Functor TreeF where
  fmap f (Node e xs) = Node e (fmap f xs)

depth = cata phi where
  phi :: Algebra TreeF Int -- TreeF Int -> Int
  phi (Node x sons) = 1 + foldr max 0 sons
```

Conclusion

Category Theory oriented Programming:

- Focus on the type and operators
- Extreme generalisation
- Better modularity
- Better control through properties of types